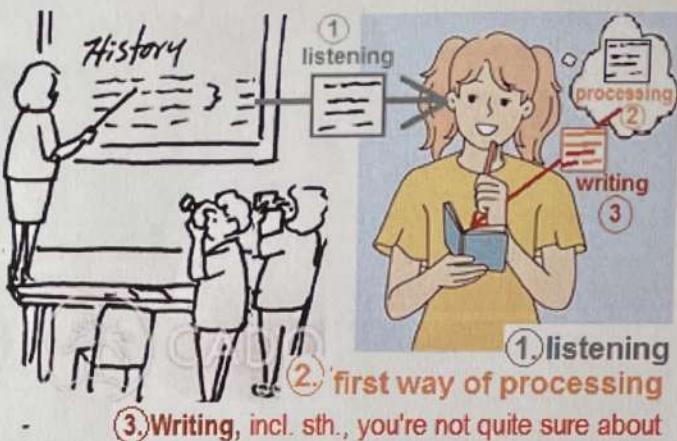
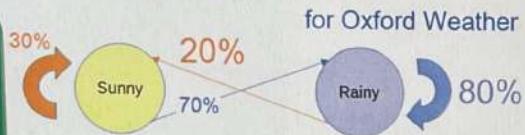
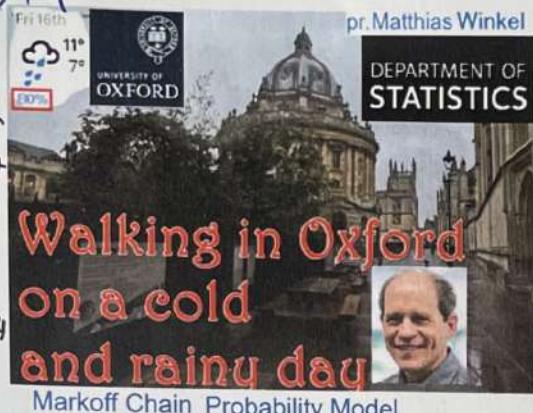


+0.1 to final Grade
+0.1



CHALK + TALK ink + think

take notes on the lecture yourself

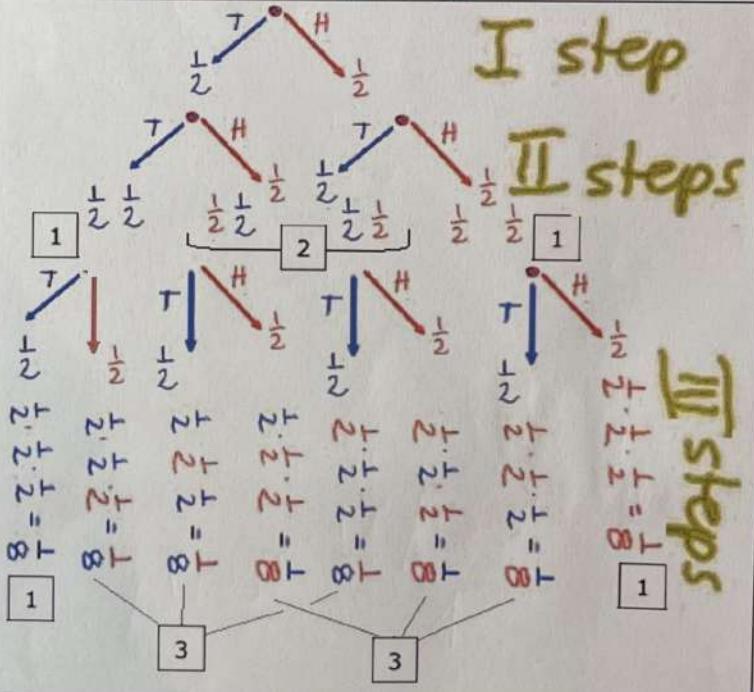
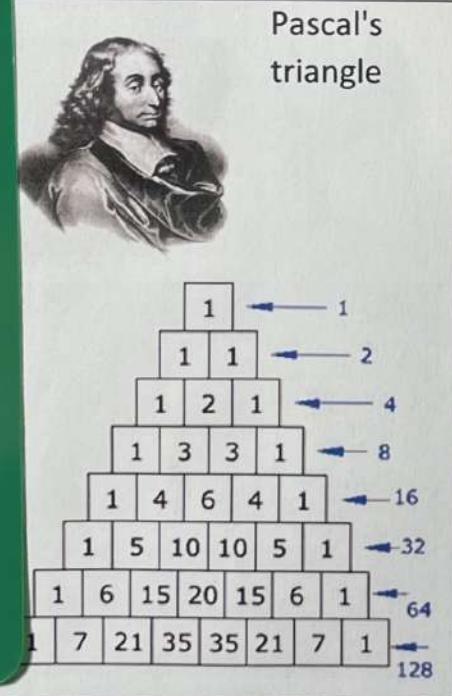
Motivation: 80% chance of rain

Let A_j be the event of rain at 9am on day j of this term, $1 \leq j \leq n$

School = formalism => $E = \{A_1, A_2, \dots, A_n\}$

Suppose the events A_i each have probability P , independently.

Oxford				
Tue 13th	Wed 14th	Thu 15th	Fri 16th	
70%	10° 9°	13° 10°	13° 8°	11° 7°
				80%



$$\begin{aligned}
 (a+b)^0 &= 1 \\
 (a+b)^1 &= a+b \\
 (a+b)^2 &= a^2 + 2ab + b^2 \\
 (a+b)^3 &= a^3 + 3a^2b + 3ab^2 + b^3 \\
 (a+b)^4 &= a^4 + 4a^3b + 6a^2b^2 + 4ab^3 + b^4 \\
 (a+b)^5 &= a^5 + 5a^4b + 10a^3b^2 + 10a^2b^3 + 5ab^4 + b^5
 \end{aligned}$$

+ 0.1
GALOIS
6!

+ 0.1
 $m_1 m_2$

w1

+ 0.1 Pr ABBA

+0.5 to final Grade

Resume of Lecture by Pr. Bob Gallagher from MIT

Massachusetts Institute of Technology (MIT)

George Boole (1815-1864) developed Boolean logic

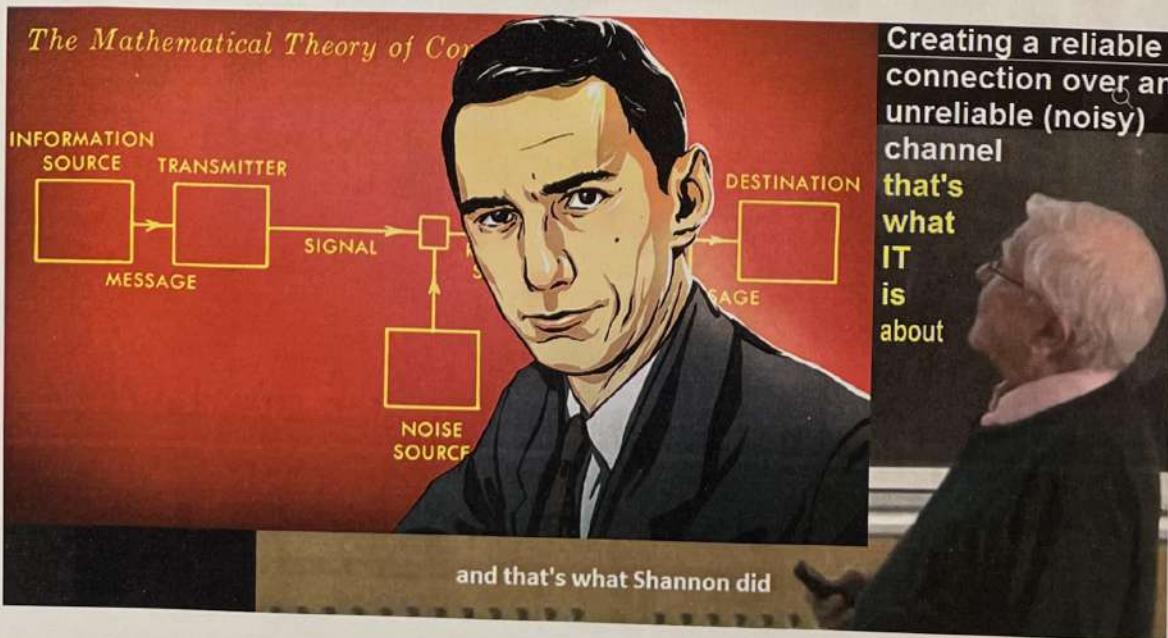
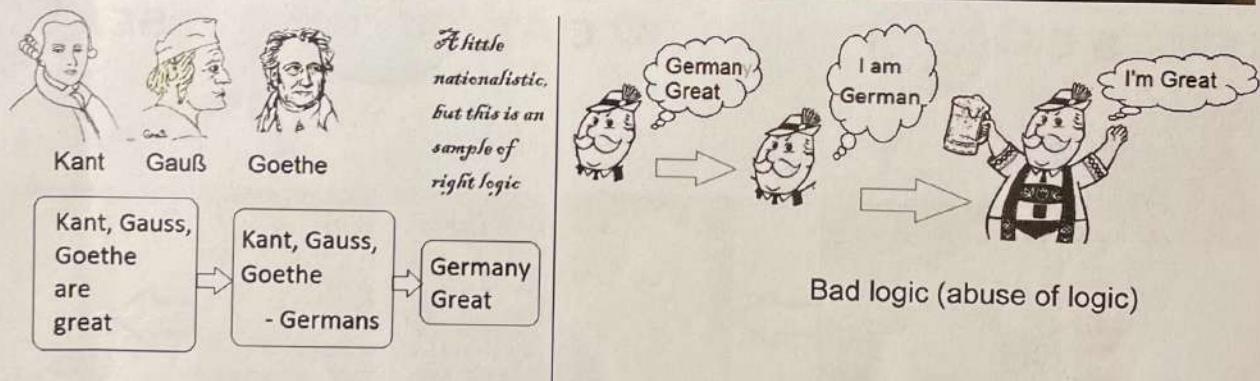
The principles of logical thinking have been understood (and occasionally used) since the Hellenic era.

Boole's contribution was to show how to systemize these principles and express them in equations (called Boolean logic or Boolean algebra).

Claude Shannon (1916-2001) showed how to use Boolean algebra as the basis for switching technology. This contribution systemized logical thinking for computer and communication systems, both for the design and programming of the systems and their applications.

Logic continues to be abused in politics, religion and most non-scientific areas

Logic continues to be abused in politics, religion, and most non-scientific areas.



(+0.2) to final grade



MIT

Massachusetts Institute
of Technology (MIT)



Lecture by Pr. Bob Gallagher

Boole (1815-1864) & Shannon (1916-2001)



9 10

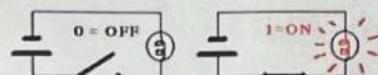
9.5

10

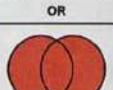
Project 0

Sapere audet
Logical addition
(disjunction)

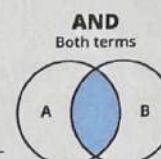
Conjunction:		
p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F



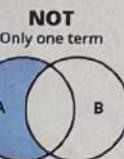
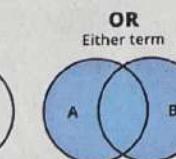
Logical disjunction



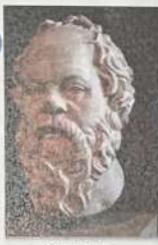
OR
Definition
 $x + y = 1$
Truth table
Logic gate



BOOLEAN LOGIC



Good logic



Socrates was
a philosopher



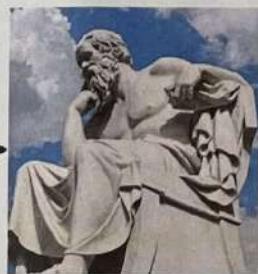
philosophers are men

$\Phi \in A$

Plato



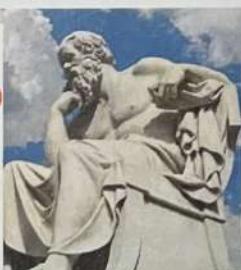
Aristotle



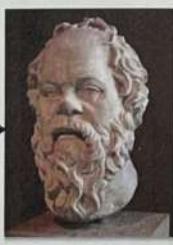
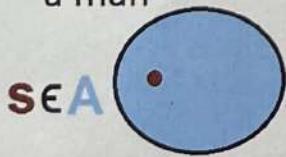
Socrates was
a man



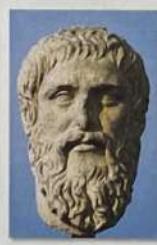
Bad logic



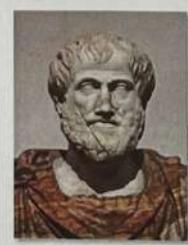
Socrates was
a man



Socrates



Plato



Aristotle

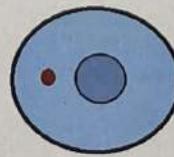


Socrates was
a philosopher



philosophers are men

$\Phi \in A$



message
source

encoder

noise

channel

decoder

message
receiver

XOR X

0 0	0
0 1	1
1 0	1
1 1	0

$$0101_2 \rightarrow 5_{10}$$

$$01111_2 \rightarrow 15$$

d: cd IT\Projects\01 ↴
↑ Any Tex

```
class MBoole
{
    public static void Main()
    {
        }
}
System.Console.WriteLine("Hi");
```

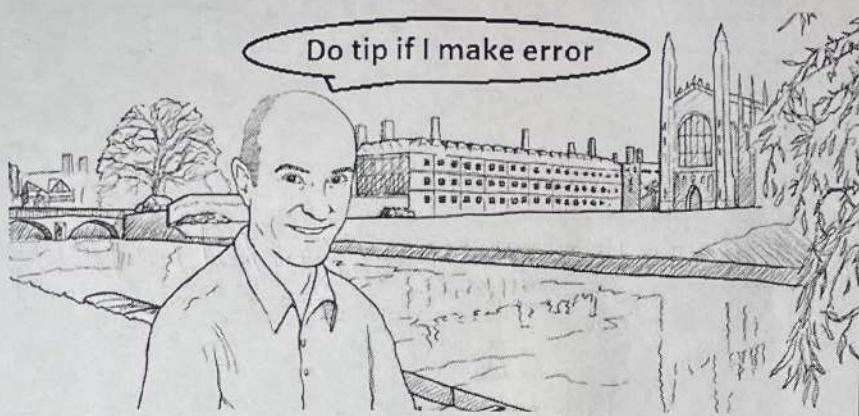
Save As

MBoole.cs

Prompt

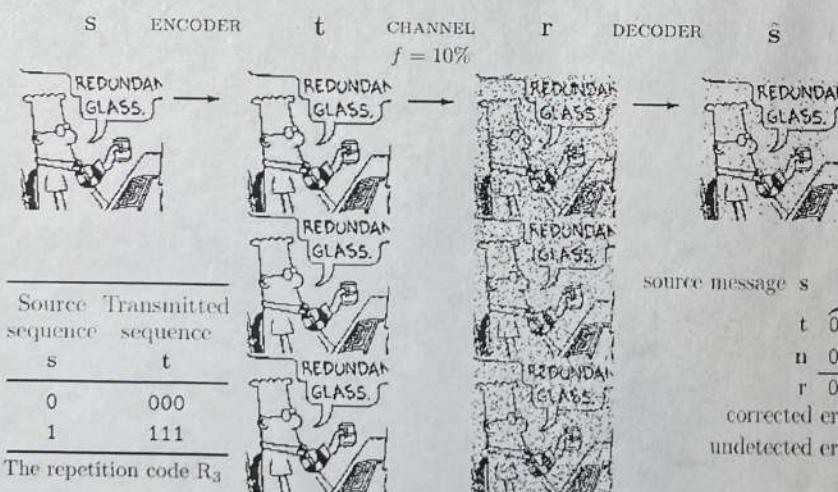
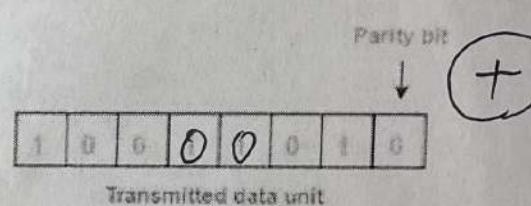
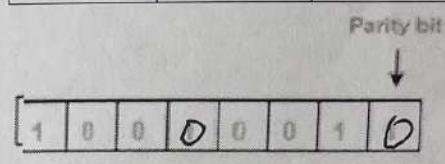
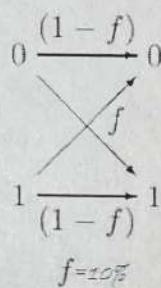
cmd d:

cd IT\Projects\01



Sir Dr. D. MacKay,
University of Cambridge
(22 April 1967 – 14 April 2016)

"I believe in clean energy,
but I also believe in mathematics"



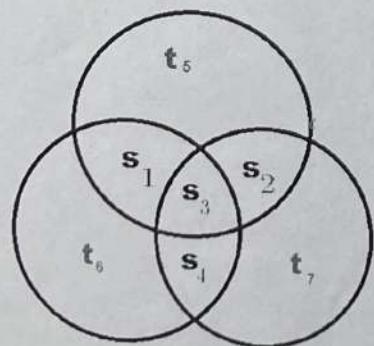
50%

ABC
DEF

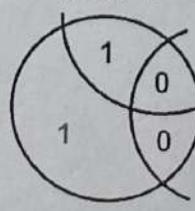
SOURCE MESSAGE	s	0	0	1	0	1	1	0
	t	000	000	111	000	111	111	000
	n	000	001	000	000	101	000	000
	r	000	001	111	000	010	111	000
	corrected errors	*						
	undetected errors		*					

7.4. Hamming code.

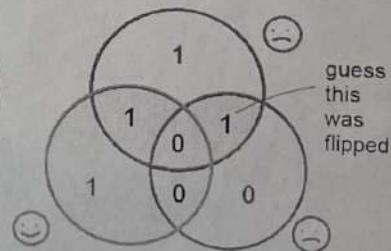
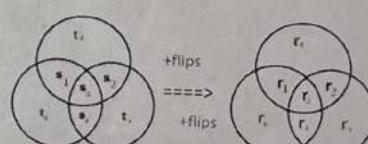
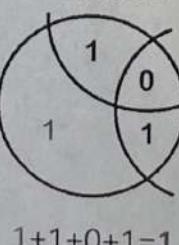
$$\frac{4}{\sum} \rightarrow \frac{7}{t}$$



satisfied

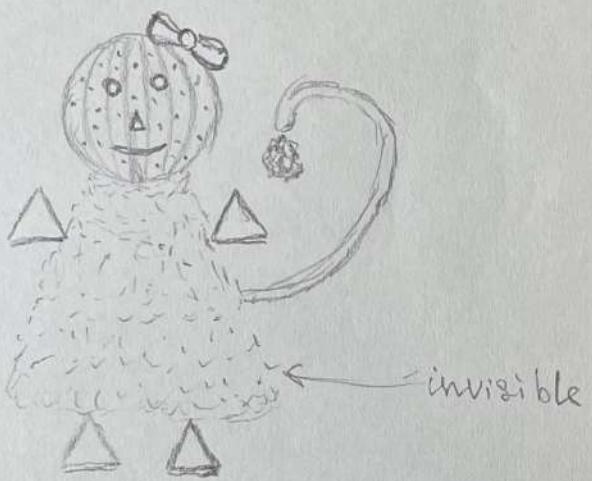


not satisfied

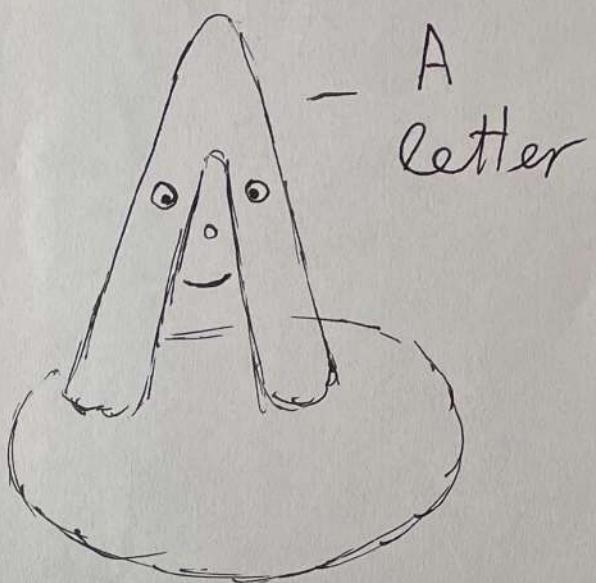


D:

+0.1 to Fin Ex



Girls
invisible

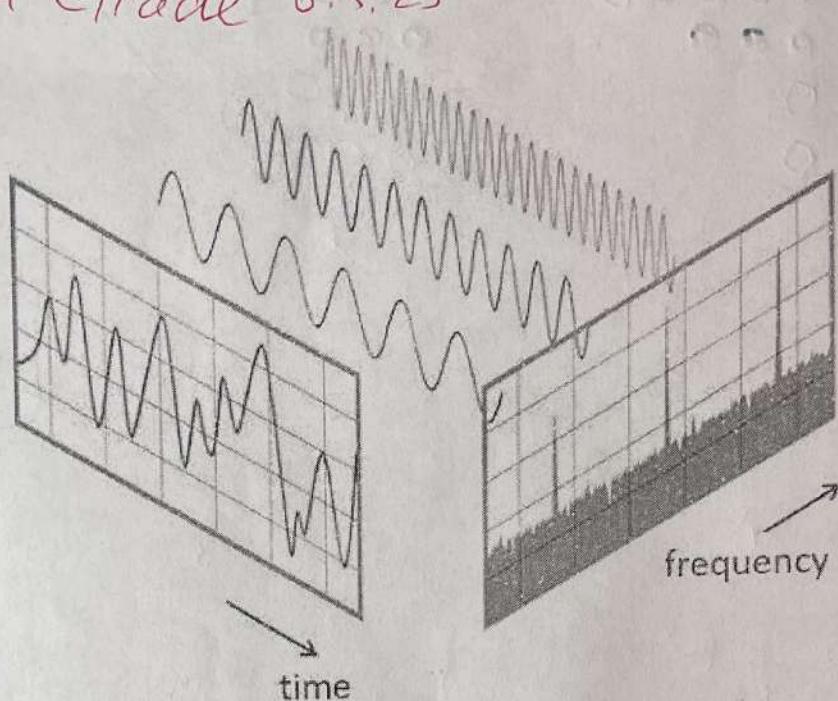


+0,7
to fin
Grade

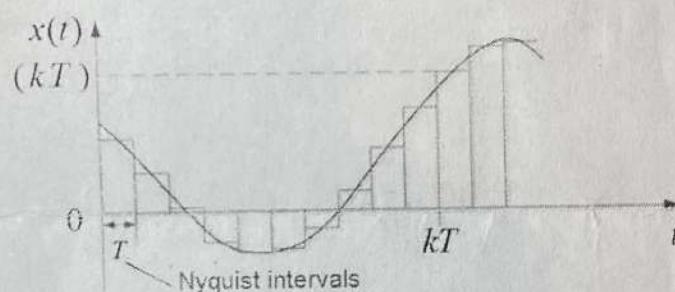
+0,1 to fin Grade 6.5.25
+0,3 to fin Grade 6.5.25

+0,1 for A
+0,1

Fourier transform

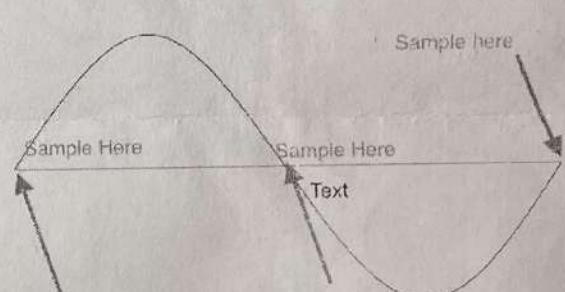


Sampling. Kotelnikov-Nyquist Theorem



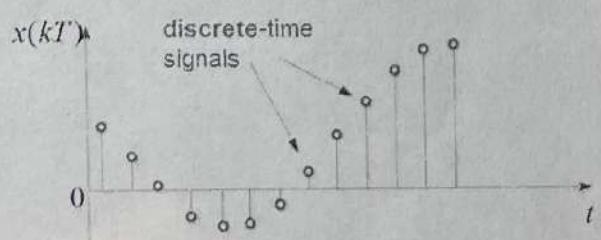
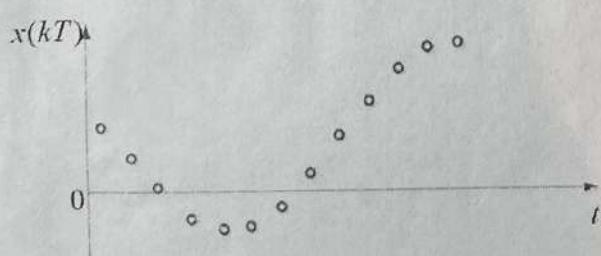
Time intervals T , through which readings s (kT) are taken, are called Nyquist intervals.

Sine with period T



Sampling at $T/2$

frequency Sample



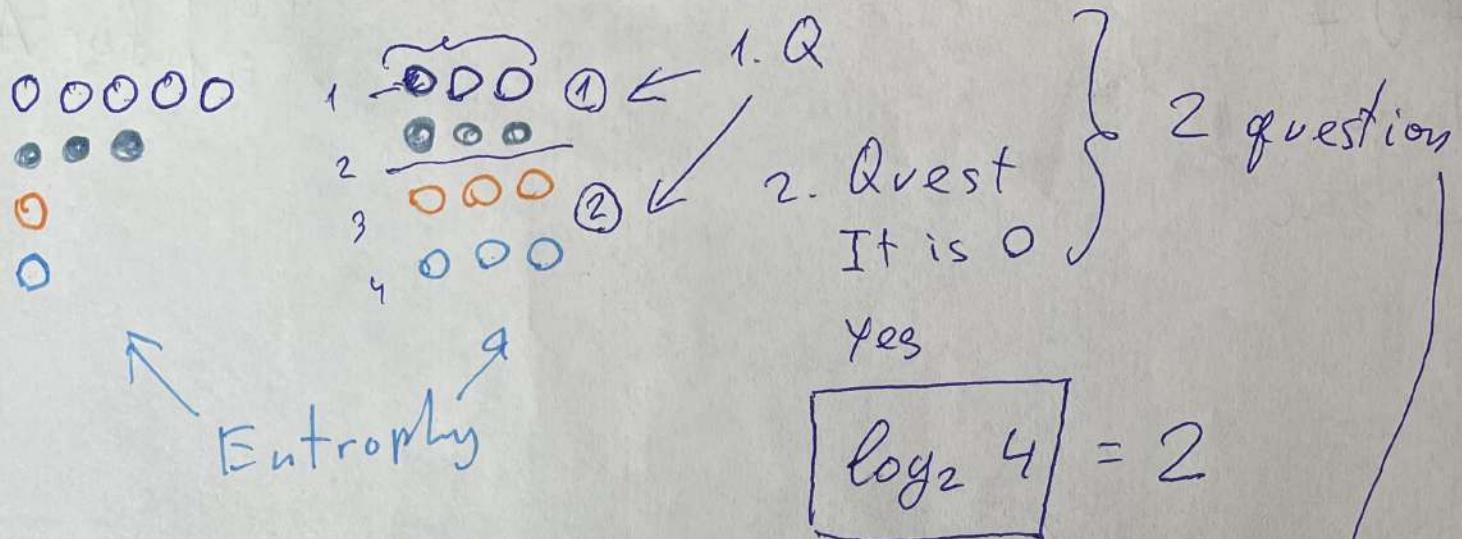
quantity of inf = $\log_2 X$

$$F_{\text{sample}} \geq 2 * F_{\text{max}} \\ (T_{\text{sample}} \leq T_{\text{min}} / 2)$$

$$\left. \begin{array}{ll} f=1 & T=1 \\ f=2 & T=\frac{1}{2} \\ f=3 & T=\frac{1}{3} \end{array} \right\} \Rightarrow T = \frac{1}{f} = f^{-1}$$

Kak nepravilnoe bie kodirovaniye b mojete,
kotorye smogut se zamenit

T - period, f - frequency (Hz)



$$1 \text{st: } H(X) = 0,5 \cdot 1 + 0,3 \cdot 2 + 0,1 \cdot 3 + 0,1 \cdot 3 = 0,5 + 0,6 + 0,3 + 0,3 = 1,7$$

$$2 \text{nd: } H(X) = \frac{3}{12} \cdot 1 + \frac{3}{12} \cdot 2 + \frac{3}{12} \cdot 3 + \frac{3}{12} \cdot 3 = 0,25 \cdot (1+2+3+3) = 0,25 \cdot 8 = 2,25$$

$$H = \sum_{i=1}^n p(x_i) \log_2 \left(\frac{1}{p(x_i)} \right) \quad - \text{Shannon Entropy}$$

$$I(x) = \log_2 \left(\frac{1}{p(x)} \right) \quad - \text{Quantifying Information}$$

2nd variant

$$\begin{aligned}
 H(X) &= 0,5 \cdot \log_2 \left(\frac{1}{0,5} \right) + 0,3 \cdot \log_2 \left(\frac{2}{0,3} \right) + 0,1 \cdot \log_2 \left(\frac{1}{0,1} \right) + 0,1 \cdot \log_2 \left(\frac{3}{0,1} \right) = \\
 &= 0,5 \cdot \log_2 2 + 0,3 \cdot \log_2 3,33 + 0,1 \cdot \log_2 10 + 0,1 \cdot \log_2 30 = \\
 &= 0,5 \cdot 1 + 0,3 \cdot 1,44 + 0,1 \cdot 3,32 + 0,1 \cdot 3,32 = \underline{1,686}
 \end{aligned}$$

Events and probabilities

- consider our "experiment", which has a set of Ω of outcomes $w \in \Omega$.

Ex. tossing a coin $\Omega = \{H, T\}$

B) throwing a die :::: :::

$$\Omega = \{(i, j), i, j \in [1, 2, 3, 4, 5, 6]\}$$

set of possible outcomes

A - subset of Ω - event

a) coin comes up tail $A = \{T\}$ - event

b) $A = \{(3, 6)\}$ - event

If $w \in \Omega$ is the outcome, we say that A occur if $w \in A$

Complement of A: \bar{A} - occur when A doesn't occur
 A^c

Set difference: $A \setminus B = A \cap B^c$

Intersection: $A \cap B$ occur if both A and B occur

Union: $A \cup B$ occur if any A or B occur

Boolean

&	00	0
	01	0
	10	0
	11	1

Union:	00	0
	01	1
	10	1
	11	1

• A and B are disjoint if $A \cap B = \emptyset$

!!

A and B can't occur together

• $p(A)$ - probability of each event A : $A = \{T\}$, $A = \{(i, j)\} \dots$

Case 1

⋮ ⋮

Ω is finite and all outcomes are equally likely

$$p(A) = \frac{|A|}{|\Omega|}$$

Example: a) coin $|\Omega| = 2$ b) dice $|\Omega| = 6$

$$|A| = 1$$

$$|A| = 6$$

$$p(A) = \frac{1}{2}$$

$$p(A) = \frac{1}{6}$$

Elementary combinatorics

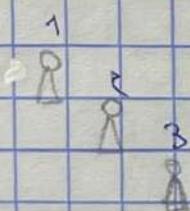
• Arrangement of distinguishable objects

• Suppose we have n distinguishable objects

1st 2nd 3rd



How many ways to order them



$$n = 3! = 6$$

For n there are n options for 1st object
then n-1 for 2nd object
inductively $n - (n-1) = n^{\text{th}}$

In all there're $n(n-1)(n-2)\dots 2 \cdot 1 = n!$ permutations

$$P(A) = \frac{m_1! \cdot m_2!}{n!}, \quad m_1 + m_2 = n \cdot 2 \cdot 3 \cdots n \\ 6 \cdot n = 2n$$

$$P(A) = \frac{(n-m)! \cdot m!}{n!}$$

Arrangements when not all objects are indistinguishable

$$\underbrace{AAA}_{\text{indistinguishable}} \quad \underbrace{BCD}_{\text{distinguishable}} \quad \} 6!$$

indistinguishable distinguishable

$$(AAA)$$

$$(A_1 A_2 A_3)$$

$$\frac{6!}{3!}$$

(number of combinations when the order
doesn't matter)

$$ABBAT = \frac{5!}{2! \cdot 2! \cdot 1!}$$

$$\frac{6!}{2!} \rightarrow \frac{6!}{2! \cdot 2!}$$

$$TALLINN = \frac{7!}{1! \cdot 1! \cdot 1! \cdot 2! \cdot 2!}$$

$$\frac{7!}{2! \cdot 2!}$$

$$m_1! \cdot m_2! \cdots m_k!.$$

$$m! \cdot (n-m)!$$

$$k=2$$

$$\frac{n!}{m! \cdot (n-m)!} \quad - \text{final formula}$$